

Dynamic Programming

Scoring matrices

BLAST

Section 3

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Outline

- Recursion
- More Dynamic Programming
- Scoring Matrices
- BLAST

Recursion

General Idea: Solve a problem by solving related sub-problems and combining the results.

Advantage: Sub-problems are easier to solve and their total run time can be much less than original problem.

Example: Merge Sort – To sort a large list, split the list in half, sort each half, and merge the two sorted lists.

Example: Factorial(x)

$$x! = x*(x-1)*(x-2)*...*1$$

$$5! = 5*4*3*2*1 = 120$$

Recursion - Factorial

Iterative Algorithm:

```
$x = 5;  
for ($n = $x-1 ; $n > 0 ; $n--) {  
  $x *= $n;  
}
```

Recursive Algorithm:

```
$x = factorial(5);  
sub factorial {  
  my($n) = $_[0];  
  if ($n <= 1) {  
    return 1; ← base case  
  } else {  
    return $n*factorial($n-1); ← recursive case  
  }  
}
```

Dynamic Programming

General Idea: Solve a problem by solving related sub-problems and combining the results.

Key Point: In recursive algorithms (e.g., Merge Sort), the sub-problems are independent. In dynamic programming the sub-problems share sub-sub-problems.

Implementation: Solve each sub-sub-problem only once and store the answers in an array.

Dynamic Programming – Step 1

Characterize the structure of the optimal solution.

Example: Sequence Alignment

Three choices at each step:

1. Align a base in X with a base in Y
2. Align a base in X with a gap
3. Align a base in Y with a gap

Goal:

Maximize the total score

Dynamic Programming – Step 2

Recursively define the value of an optimal solution.

Example: Needleman-Wunsch (Global Alignment)

$$F(i,j) = \max \{ \begin{array}{l} F(i-1, j-1) + s(x_i, y_j) , \\ F(i-1, j) - d , \\ F(i, j-1) - d \end{array} \}$$

Dynamic Programming – Step 3

Compute the optimal solution, starting from the

easiest sub-sub-problems and saving the results.

	-	A	C	A	C	T
-	0	-3	-6			
A	-3	2				
G	-6					
C						
A						
C						

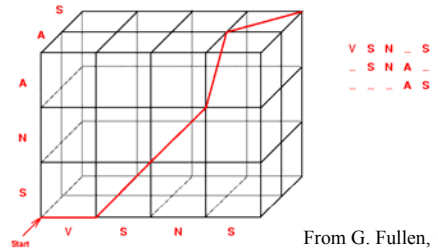
Dynamic Programming – Step 4

Construct an optimal solution from the computed information. (Perform a trace-back.)

	-	A	C	A	C	T
-	0	-3	-6	-9	-12	-15
A	-3	2	-4	-7	-10	
G	-6	-1	1	-5	-8	
C	-9	-4	1	0	-3	
A	-12	-7	-2	3	0	
C	-15	-10	-5	0	5	2

Multiple Sequence Alignment

Dynamic programming in many dimensions!!!



From G. Fullen, 1996.

Dynamic Programming Summary

Advantage: Guaranteed to find the mathematically optimal solution (highest scoring alignment). Much faster than trying all possible alignments.

Disadvantage: For large problems (long sequences or many sequences), dynamic programming can be very slow and require lots of computer memory.

Scoring Matrices

Definition: A table (matrix) of values that describe the probability of a residue (amino acid or base) pair occurring in an alignment of related sequences.

Why a matrix? Evolution tends to favor different sets of substitutions. Thus, a single probability does not accurately describe all residue pairs

BLAST : K-mer Hashing

- For each k-mer, list all locations where it occurs in the database.
- Match the query to k-mers (or “neighborhood words” above a threshold) in the list. This forms seeds.

The BLAST Search Algorithm



BLAST : Extending Seeds

- Connect “close” seeds into single long sequences.
- Extend seeds until the cumulative score drops.
- High-scoring Segment Pair = locally maximal segment.

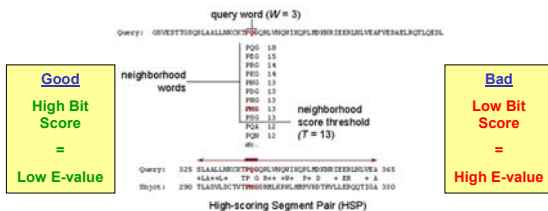
The BLAST Search Algorithm



BLAST : Evaluating HSPs

- Bit Score: independent of the scoring matrix used.
- E-value: “Expected number of High-Scoring Segment Pairs”. The smaller the number, the better the alignment.
- E-values of 0.1 or 0.05 are typically used as cutoffs.

The BLAST Search Algorithm



Next Week

- Microarrays
- Sequencing

Acknowledgement / References

This handout includes material written by Suzanne Komili, Yonatan Grad, Doug Selinger, and Zhou Zhu.

Mount, *Bioinformatics – Sequence and Genome Analysis*; Durbin et al., *Biological Sequence Analysis*; Cormen et al., *Introduction to Algorithms*

<http://www.ncbi.nlm.nih.gov/BLAST>